



Security Engineering Fall 2015

Lecture 15 – DB Security
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(Most material courtesy of D. Gollmann)





Lecture Outline

DB security

- General threats
- Relational databases countermeasures
- Computation on encrypted data

Statistical DB security

- Inferential threats
- Statistical databases countermeasures

New challenges





Database Security: Statistics

- 621 out of 47000+ security incidents lead to data disclosures and to 44 million+ compromised records
- Main threats
 - Malware 40%
 - Hacking 52%
 - Social 29%
 - Misuse 13%
 - Physical 35%
 - Error 2%
 - Environmental 0%





Database Security

- A database is a collection of data
- DBMS organizes the data and gives users the means to retrieve information
- Database Security:
 - protection of sensitive data and mechanisms that allow users to retrieve information in a controlled manner
 - Provide controlled, protected access to the contents of a database as well preserve the integrity, consistency and overall quality of the data





Design Requirements

• Precision:

 protect sensitive information while revealing as much nonsensitive information as possible

Internal consistency:

- the entries in the database obey some prescribed rules
 - E.g., stock levels cannot fall below zero.

• External consistency:

- The entries in the database are correct
 - E.g., stock levels given in the database match stock levels in the warehouse
- DBMS alone cannot keep the database in a consistent state
 → need organizational controls!
- This property is also called accuracy





Top Ten Database Security Threats

(lack of) organzational measures

- 1. Excessive and Unused Privileges
- 2. Abuse Privileges
- 8. Unmanaged Sensitive Data
- 10. Limited Security Expertise and Education

(presence of) software problems

- 3. SQL Injection
- 4. Malware

5. A bit of both

- 5. Weak Audit Trail
- 6. Storage Media Exposure
- 7. Exploitation of (unpatched) Vulnerabilities
- 9. Denial of Service





T1: Excessive and Unused Privileges

- When someone is granted database privileges that exceed the requirements of their job function, these privileges can be abused
- Access control mechanisms for job roles have not been well defined or maintained
 - users may be granted generic or default access privileges that far exceed their specific job requirements

Example

 a bank employee whose job requires the ability to change only account holder contact information may take advantage of excessive database privileges and increase the account balance of a colleague's savings account





T2: Abuse Privileges

 Users will abuse legitimate database privileges for unauthorized purposes

Example

- Consider an internal healthcare application used to view individual patient records via a custom Web interface
- The Web application normally limits users to viewing an individual patient's healthcare history
- However, a rogue user might be able to circumvent these restrictions and copy electronic healthcare records on his laptop





T8: Unmanaged Sensitive Data

- Many companies struggle to maintain an accurate inventory of their databases
 - Forgotten databases may contain sensitive information
 - New databases can emerge
- Sensitive data in these databases will be exposed to threats if the required controls and permissions are not implemented
 - You cannot control what you don't "know" it exists



Limited Security Expertise and Education

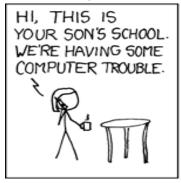
- Internal security controls are not keeping pace with data growth and many organizations are ill-equipped to deal with a security breach
- Due to the lack of expertise required to implement security controls, policies, and training
- According to PWC's 2012 Information Security Breaches Survey
 - 75% of the organizations surveyed experienced staffrelated breaches when a security policy was poorly understood
 - 54% of small businesses did not have a program for educating their staff about security risks





T3: SQL Injection

- Already discussed for application security
 - SQL injection involves inserting (or "injecting") unauthorized or malicious database statements into a vulnerable SQL data channel such as a Web application or stored procedure
 - If these injected statements are executed by the database, critical data stores can be viewed, copied, and altered
- Normally this happens by exploiting the polyglottism of the web application on top of the DB
 - http://xkcd.com/327













T4: Malware

- Cybercriminals, state-sponsored hackers, and spies use advanced attacks to penetrate organizations
 - spear phishing emails and malware
- Legitimate users become a conduit for these groups to access networks and sensitive data
 - Users are unaware that malware has infected their device





T5: Weak Audit Trail

- Failure to collect detailed audit records of database activity
 - Organizations with weak (or sometimes non-existent) database audit mechanisms will increasingly find that they are at odds with industry and government regulatory requirements
- Many enterprises will turn to native audit tools provided by their database vendors or rely on ad-hoc and manual solutions
 - These approaches do not record details necessary to support auditing, attack detection, and forensics.
- Finally, users with administrative access to the database, either legitimately or maliciously obtained, can turn off native database auditing to hide fraudulent activity
 - Audit duties should ideally be separate from both database administrators and the database server platform to ensure strong separation of duties policies





T6: Storage Media Exposure

- Backup storage media is often completely unprotected from attack
 - Numerous security breaches have involved the theft of database backup disks and tapes
- Failure to audit and monitor the activities of administrators who have low-level access to sensitive information can put your data at risk
 - Taking the appropriate measures to protect backup copies of sensitive data
 - Monitor your most highly privileged users is not only a data security best practice, but also mandated by many regulations





T7: Exploitation of Vulnerabilities

- It is common to find vulnerable and un-patched databases, or discover databases that still have default accounts and configuration parameters
 - Attackers know how to exploit these vulnerabilities to launch attacks against your organization
- Maintenance is hard
 - Organizations often struggle to stay on-top of maintaining database configurations even when patches are available
 - It generally takes organizations months to patch databases once a patch is available
 - Sometimes licensing is the issue, sometimes interoperability with legacy software





T9: Denial of Service

- DoS conditions can be created via many techniques
- Most common technique
 - overload server resources such as memory and CPU by flooding the network with database queries that ultimately cause the server to crash
- Motivations behind DoS attacks
 - often linked to extortion scams in which a remote attacker will repeatedly crash servers until the victim meets their demands
- Software or organizational decision
 - Not enough "power" to meet excess demand





Relational Databases

- A relational database is a database that is perceived by its users as a collection of tables
- A relation R is a subset of $D_1 \times \cdots \times D_n$ where D_1, \ldots, D_n are the domains on n attributes.
- The elements in the relation are n-tuples (v_1, \ldots, v_n) with $v_i \in D_i$; the value of the i-th attribute has to be an element from D_i .
- Elements in a tuple are often called fields
- A special null value indicates that a field does not contain any value





Example

Name	Day	Flight	Status
Alice	Mon	GR123	private
Bob	Mon	YL011	business
Bob	Wed	BX201	
Carol	Tue	BX201	business
Alice	Thu	FL9700	business

Flight	Destination	Departs	Days
GR123	THU	7: 55	14
YL011	ATL	8:10	12345 - 7
BX201	SLA	9: 20	1 - 3 - 5
FL9700	SLA	14: 00	- 2 - 4 - 6 -
GR127	THU	14: 55	-2 - 5-





SQL

Structured Query Language (SQL):

 standard language for describing how information in a relational database can be retrieved and updated.

SQL operations:

- SELECT: retrieves data from a relation.
- UPDATE: update fields in a relation.
- DELETE: deletes tuples from a relation.
- INSERT: adds tuples to a relation.





Examples

SELECT Name, Status

FROM Diary

WHERE Day = 'Mon'

UPDATE Diary

SET Status = private

WHERE Day = 'Sun'

DELETE FROM Diary

WHERE Name = 'Alice'

INSERT INTO Flights (Flight, Destination, Days)

VALUES ('GR005', 'GOH', '12-45-')





Types of Relations

• Base relations (real relations):

- named, autonomous relations; exist in their own right, are not derived from other relations,
- have 'their own' stored data

Views:

- named, derived relations, defined in terms of other named relations;
- no stored data of their own.

Snapshots:

- named, derived relations, defined in terms of other named relations;
- have stored data of their own.

Query results:

- may or may not have a name;
- no persistent existence in the database per se





Database Keys

- Tuples in a relation must be uniquely identifiable
- A primary key K of a relation R must be:
 - Unique: at any time, no tuples of R have the same value for K;
 - Minimal: if K is composite, no component of K can be omitted without destroying uniqueness.
- Every relation must have a primary key
- A primary key of one relation that is an attribute in some other relation is a foreign key in that relation





Integrity Rules

Entity Integrity Rule:

 no component of primary key of a base relation is allowed to accept nulls

Referential Integrity Rule:

DB must not contain unmatched foreign key values

Application specific integrity rules:

- Field checks: to prevent errors on data entry
- Scope checks
- Consistency checks





Access Control

- Discretionary access control using privileges and views, based on:
 - users: authenticated during logon;
 - actions: include SELECT, UPDATE, DELETE, and INSERT;
 - objects: tables, views, columns (attributes) of tables and views;
- Users invoke actions on objects
- DBMS decides whether to permit the requested action
- When an object is created, it is assigned an owner; initially only the owner has access to the object; other users have to be issued with a privilege:
 - (grantor, grantee, object, action, grantable)





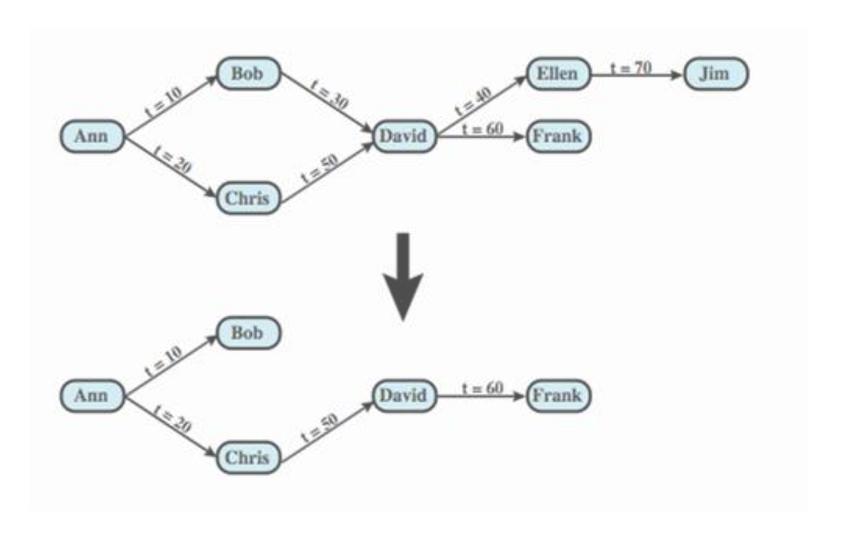
Granting & Revoking Privileges

- Privileges managed with GRANT and REVOKE
 - GRANT SELECT, UPDATE (Day, Flight)
 - ON TABLE Diary
 - TO Art, Zoe
- Selective revocation of privileges:
 - REVOKE UPDATE
 - ON TABLE Diary
 - FROM Art
- Right to delegate privileges given through GRANT option:
 - GRANT SELECT
 - ON TABLE Diary
 - TO Art
 - WITH GRANT OPTION





Granting and Revoking Privileges







Role-Based Access Control

- Role-based access control work well for DBMS
 - eases admin burden, improves security
- categories of database users:
 - application owner
 - end user
 - administrator
- DB RBAC must manage roles and their users
 - cf. RBAC on Microsoft's SQL Server
- It can be implemented using GRANT and REVOKE





Row Level Access Control

- Views: derived relations, created by
 - CREATE VIEW view_name [(column [, column] ...)]
 - AS subquery
 - [WITH CHECK OPTION]
- Many security policies better expressed by privileges on views than by privileges on base relations
- Access conditions described through subquery in the view definition
 - CREATE VIEW business_trips AS
 - SELECT * FROM Diary
 - WHERE Status = `business'
 - WITH CHECK OPTION;





Advantages

- Views are flexible and allow access control policies to be defined at a level of description that is close to the application requirements
- Views can enforce context-dependent and data-dependent security policies.
- Views can implement controlled invocation
- Secure views can replace security labels
- Data can be easily reclassified





Examples

Application-level access control

```
CREATE VIEW Top of the Class AS
 SELECT
                     displays students whose grade
 FROM Students
                     average is higher than that of the person using the view
 WHERE Grade > (SE
                  FRom Bludents
                  WHERE Name
                                 display journeys booked
                                 by the customer using the
CREATE VIEW My_Journeys
      SELECT * FROM Diary
       WHERE Customer = current user());
```





CHECK Option

- INSERT and UPDATE can interfere with view-based access control
- Views may not be updatable because they do not contain the information that is needed to maintain the integrity of the corresponding base relation
 - E.g., a view that does not contain the primary key of an underlying base relation cannot be used for updates
- Blind writes:
 - updates that overwrite an existing entry
 - For views defined WITH CHECK OPTION, UPDATE and INSERT can only write entries to the database that meet the definition of the view
- Blind writes possible if CHECK option is omitted





CHECK Option: Example

- UPDATE business_trips AS
- CREATE WEW business_trips AS

WHIERE NEROM DATINE' AND Day = 'Thu'

WHERE Status = 'business'

WITH CHECK OPTION

Name	Day	Flight	Status
Bob	Mon	YL011	business
Carol	Tue	BX201	business
Alice	Thu	FL9700	business





Database Encryption

• Databases are a valuable information resource

 protected by multiple layers of security: firewalls, authentication, O/S access control systems, DB access control systems, and database encryption

Encryption strategies

- entire database very inflexible and inefficient
- individual fields simple but inflexible
- records (rows) or columns (attributes) best
 - also need attribute indexes to help data retrieval

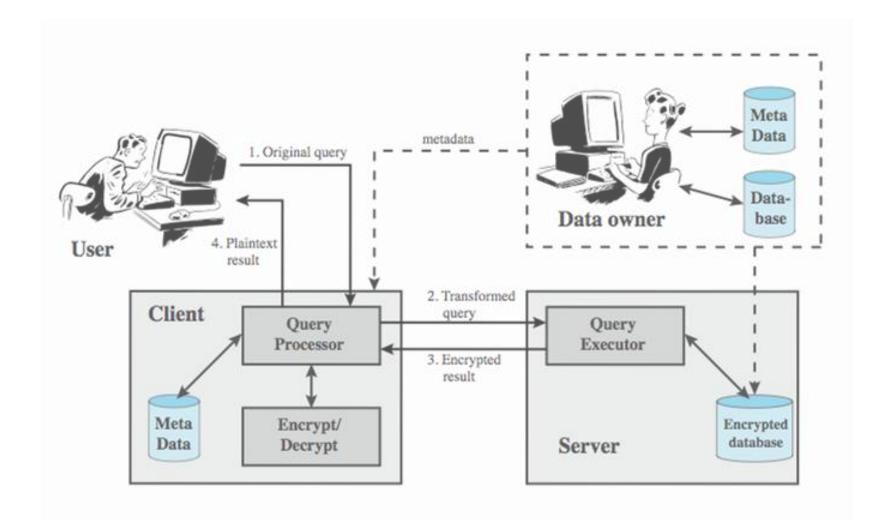
Issues

- Key Management
- Query Execution





Database Encryption







Dtabase encryption functionalities

• Fragments of SQL can be captured by different kind of encryptions

- Order-preserving encryption preserves the order of plaintexts, i.e.
 - $x \le y \Rightarrow E_{OP}(x) \le E_{OP}(y)$
- Deterministic encryption preserves the equality of plaintexts, i.e.
 - $x=y \Rightarrow E_{DET}(x) = E_{DET}(y)$
- In (additively) homomorphic encryption multiplication of ciphertexts (modulo a key-dependent constant) maps to addition of the plaintexts,
 - $D_{HOM}(E_{HOM}(x)*E_{HOM}(y)) = x+y$

Unfortunately not all types of encryptions can be combined

- SELECT x FROM T GROUP BY y HAVING SUM(z) > 100
- Requires complex constructions data is layered in onions of encryption

Not all encryption schemes equally secure (why?)

- OP ≤ DET ≤ HOM
- However not everything needs to be encrypted with less secure scheme (but makes DB very dependent on type of query, and cannot change later)

• Performance is still painful

- SAP performance Data (GIS 2014)
 - Exact Search 1.2x, Equi-Join 1.5x, Group&Sum 11.7x, OrderbySum 15.4x

Harder with "standard" access control on the DB for multiple users

- Different users have access to different fragments of the DB which therefore must use different keys
- SAP implementation by F. Kershbaum is +37% over single encrypted users





Statistical Database Security

Statistical database

 information retrieved by means of statistical (aggregate) queries on attributes (columns) of a table.

Aggregate functions in SQL

- COUNT: the number of values in a column,
- SUM: the sum of the values in a column,
- AVG: the average of the values in a column,
- MAX: the largest value in a column,
- MIN: the smallest value in a column.

Query predicate of a statistical query

specifies the tuples used for computing the aggregate value

Query set

tuples matching the query predicate





Threats

- Key difference with normal databases: user is authorized to ask stuff but not too much
- Aggregation
 - sensitivity level of an aggregate computed over a group of values may differ from the sensitivity levels of the individual elements;
 - e.g., an aggregate may be sensitive information derived from a collection of less sensitive business data

Inference

- Derivation of sensitive information from non-sensitive data
- Direct Attack:
 - aggregate computed over a small sample so that information about individual data items is leaked
- Indirect Attack:
 - combine information relating to several aggregates
- Tracker Attack:
 - a particularly effective type of indirect attack





Direct Attack

Name	Sex	Program	Units	Grade Ave.
Alma	F	MBA	8	63
Bill	М	CS	15	58
Carol	F	CS	16	70
Don	М	MIS	22	75
Errol	М	CS	8	66
Flora	F	MIS	16	81
Gala	F	MBA	23	68
Homer	М	CS	7	50
Igor	М	MIS	21	70





Direct Attack

Assume we know Carol is a female CS student

Q1 : SELECT COUNT(*)

FROM Students

WHERE Sex = 'F' AND Returns count 1

Q2 : SELECT AVG (Grade Ave.)

FROM Stu

Returns 70: average
for a single student

m = 'CS'





Tracker Attack

Q3 : SELECT COUNT(*)

FROM Students

WHERE Programme = 'CS'

Q4 : SELECT COUNT(*)

FROM Students

WHERE Programme = 'CS' AND Sex = 'M'

Q5 : SELECT AVG (Grade Ave.)

FROM Students

WHERE Program = 'CS'

Returns count 4

Returns count 3

Returns average 58

Returns average 61

Q6 : SELECT AVG (Grade Ave.)

FROM Students

WHERE Pro Carol's grade average:

 $4 \cdot 61 - 3 \cdot 58 = 70$





General Attacks

- Individual tracker for a given tuple
 - Query R that allows to derive information about our tuple r
- General tracker
 - Query T used to find the answer to any inadmissible query.
- How to combine them?
 - T chosen so that the query set and its complement are large enough for the query to be permitted
 - Make two queries to the database with the predicates R \vee T and R \vee \neg T;
 - the target r is the only tuple used by both queries
- 'Add' the two results and 'subtract' the result of a query over the entire database
- Only the target is left





General Tracker

Q7 : SELECT SUM(Units)

FROM Students

WHERE Name = 'Carol' OR Program = 'MIS'

Q8 : SELECT SUM(Units)

FROM Students

WHERE Name = 'Carol' OR NOT (Program = 'MIS')

Q9 : SELECT SUM(Units)

FROM Students

Returns sum 136

Returns sum 75

Returns sum 77

Carol has passed (75 + 77) - 136 = 16 units





Countermeasures

- Suppress obviously sensitive information
- Disguise the data
 - Randomly modify entries in the database so that an individual query will give a wrong result although the statistical queries still would be correct
 - R. Agrawal, R. Srikant: Privacy-Preserving Data Mining. SIGMOD 2000.
 - Add small random perturbations to query result so that the value returned is close to the real value but not quite correct
 - Drawback: reduced precision and usability
- Better design of the database schema
- Track what the user knows
 - user actions recorded in an audit log, a query analysis checks for suspicious sequences of queries





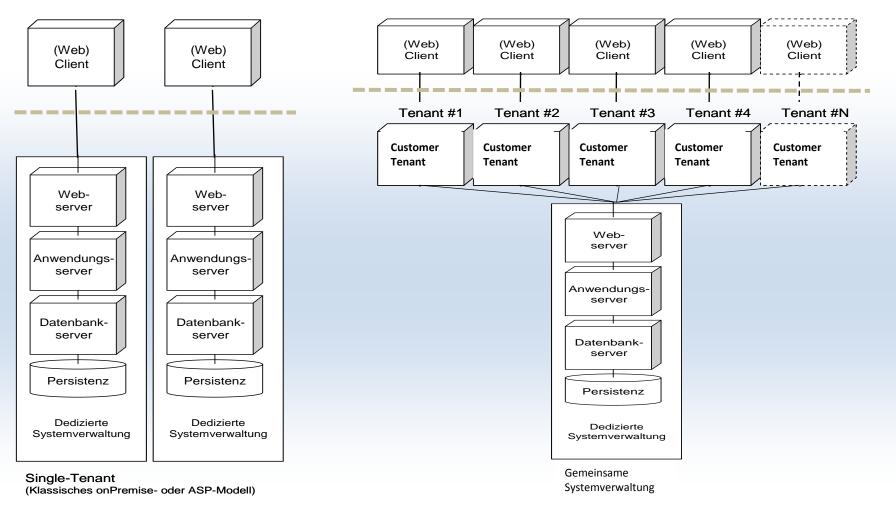
Life is complicated

- Protecting access to services requires protecting all layers
 - Web server
 - Application Server
 - DBMS Server
 - OS System + Hw Infrastructure
- They used to be isolated into a silo for different customers
 - The "Apache" model
 - Things change with cloud models





From ASP to Multi-Tenancy



Single Tenancy (classical on-premise or ASP model)

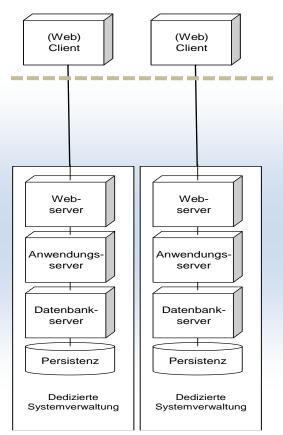
Multi Tenancy

Source: SAP

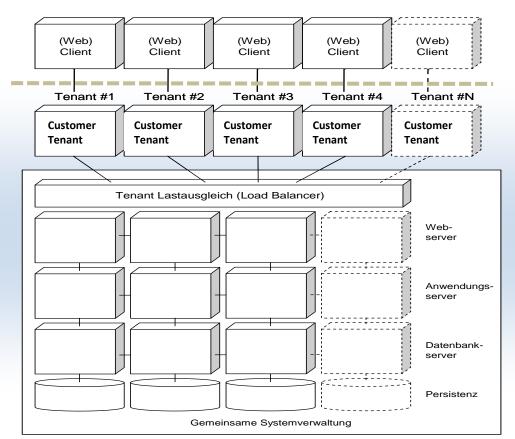




Efficient & Sclabale Multi-Tenancy



Single Tenancy (classical on-premise or ASP model)



Multi Tenancy

Source: SAP





Additional readings

- Chapter 9. D. Gollmann Computer Security
- Chapter 5 W. Stallings & L. Brown Computer Security -Principles and Practices
- Animated Database Courseware
 - http://adbc.kennesaw.edu/
- Top 10 Database Security Threats
 - http://www.imperva.com/docs/WP TopTen Database Threats.
 pdf
- Searching on Encrypted data
 - http://www.fkerschbaum.org/seed.html
- Search Google for DataCenter Security
 - http://www.youtube.com/watch?v=1SCZzgfdTBo